

Lecture 16 Semidirect products

First some terminology. Let G be a group.

An automorphism of G is an isomorphism $\varphi: G \rightarrow G$.
The identity function $I: G \rightarrow G$ is an automorphism, but there may be others. For instance:

Let $g \in G$. Then $c_g: G \rightarrow G$, $c_g(x) = gxg^{-1}$ is called conjugation by g .

Lemma: c_g is an automorphism of G .

Proof $c_g(xy) = gxyg^{-1} = gxg^{-1}gyg^{-1} = c_g(x)c_g(y)$

The inverse of c_g is $c_{g^{-1}}$:
 $c_g(c_{g^{-1}}(x)) = g(g^{-1}xg)g^{-1} = x$
 $c_{g^{-1}}(c_g(x)) = g^{-1}(gxg^{-1})g = x.$

so c_g is a bijective homomorphism. \square

We denote by $\text{Aut}(G)$ the set of all automorphisms of G .
 $\text{Aut}(G)$ is a subset of $\text{Sym}(G) = \{f: G \rightarrow G \mid f \text{ is bijective}\}$,
and in fact it is a subgroup. (The composition of automorphisms is an automorphism, and the inverse of an automorphism is an automorphism - check this.)

Proposition The function $\alpha: G \rightarrow \text{Aut}(G)$, $\alpha(g) = c_g$ is a homomorphism.

Proof $\alpha(g_1 g_2)(x) = c_{g_1 g_2}(x) = (g_1 g_2)x(g_1 g_2)^{-1}$
 $= g_1 g_2 x g_2^{-1} g_1^{-1}$

$\alpha(g_1) \circ \alpha(g_2)(x) = c_{g_1}(c_{g_2}(x)) = g_1(g_2 x g_2^{-1})g_1^{-1}$
 $= g_1 g_2 x g_2^{-1} g_1^{-1}$ same.

Proposition A subgroup $N \subseteq G$ is normal if and only if
 $c_g(N) = N$ for all $g \in G$

Proof Suppose $c_g(N) = N$. Take $n \in N, g \in G$.
 then $gng^{-1} \in c_g(N) = N$ so N is normal.

Suppose N normal. Then $c_g(N) = \{c_g(n) \mid n \in N\} = \{gng^{-1} \mid n \in N\} \subseteq N$
 (for any $g \in G$). Then $c_{g^{-1}}(N) \subseteq N$ as well, and

$$N = c_g(c_{g^{-1}}(N)) \subseteq c_g(N). \text{ thus } N = c_g(N)$$

Corollary If $N \triangleleft G$, there is a homomorphism
 $\alpha: G \rightarrow \text{Aut}(N)$, $\alpha(g) = (c_g \text{ restricted to } N)$.

Semidirect products Like direct products, we start with
 two groups A, N (not necessarily subgroups of some third group.)
 We also choose a homomorphism

$$\alpha: A \rightarrow \text{Aut}(N)$$

We then consider the set $N \times A = \{(n, a) \mid n \in N, a \in A\}$
 and we define a binary operation.

$$(n, a)(n', a') = (n \cdot \alpha(a)(n'), a a')$$

NB. α is a function $A \rightarrow \text{Aut}(N)$ so $\alpha(a) \in \text{Aut}(N)$
 so $\alpha(a)$ is a function $N \rightarrow N$, and
 $\alpha(a)(n')$ is an element of N .

A somewhat nicer notation is to write α_a in place of $\alpha(a)$.
 So then the binary operation is written:

$$(n, a)(n', a') = (n \alpha_a(n'), aa')$$

Proposition This operation makes the set $N \times A$ into a group.
 We denote it by $N \rtimes_{\alpha} A$.

Proof Associative: $((n_1, a_1)(n_2, a_2))(n_3, a_3) = (n_1 \alpha_{a_1}(n_2), a_1 a_2)(n_3, a_3)$
 $= (n_1 \alpha_{a_1}(n_2) \alpha_{a_1 a_2}(n_3), a_1 a_2 a_3)$
 $= (n_1 \alpha_{a_1}(n_2) \alpha_{a_1}(\alpha_{a_2}(n_3)), a_1 a_2 a_3)$
 $= (n_1 \alpha_{a_1}(n_2 \alpha_{a_2}(n_3)), a_1 a_2 a_3)$
 $= (n_1, a_1)(n_2 \alpha_{a_2}(n_3), a_2 a_3)$
 $= (n_1, a_1)((n_2, a_2)(n_3, a_3))$

\downarrow since $\alpha: A \rightarrow \text{Aut}(N)$ is homomorphism
 \downarrow since $\alpha_a: N \rightarrow N$ is a homomorphism.

Identity: $(e, e)(n, a) = (e \alpha_e(n), ea) = (en, ea) = (n, a)$
 $(n, a)(e, e) = (n \alpha_a(e), ae) = (ne, ae) = (n, a)$

Inverses $(n, a)(\alpha_{a^{-1}}(n^{-1}), a^{-1}) = (n \alpha_a \alpha_{a^{-1}}(n^{-1}), aa^{-1})$
 $= (n \alpha_{aa^{-1}}(n^{-1}), e) = (n \alpha_e(n^{-1}), e) = (nn^{-1}, e) = (e, e)$
 similarly $(\alpha_{a^{-1}}(n^{-1}), a^{-1})(n, a) = (e, e)$ \square

How to remember the formula:

$$\begin{array}{ccc} (n, a)(n', a') & & \\ \downarrow & \swarrow & \downarrow \\ (n \alpha_a(n'), a a') & & \end{array}$$

When the a passes over the n' ,
we apply α_a to n' .

The direct product $N \rtimes A$ is the special case where
 $\alpha: A \rightarrow \text{Aut}(N)$ is trivial: $\alpha_a = (\text{Identity } N \rightarrow N)$ for all a .

Properties of $N \rtimes_{\alpha} A$: Let $\tilde{N} = \{(n, e) \mid n \in N\}$
 $\tilde{A} = \{(e, a) \mid a \in A\}$

Then $\tilde{N} \cong N$, $\tilde{A} \cong A$, \tilde{N} and \tilde{A} are subgroups of $N \rtimes_{\alpha} A$,
and \tilde{N} is a normal subgroup of $N \rtimes_{\alpha} A$,
while \tilde{A} is not necessarily normal.
 $\tilde{N} \cap \tilde{A} = \{(e, e)\}$, and $\tilde{N}\tilde{A} = N \rtimes_{\alpha} A$.

Most interesting bit is that \tilde{N} is normal: indeed

$$\begin{aligned} (m, b)(n, e)(\alpha_{b^{-1}}(m^{-1}), b^{-1}) \\ = (m, b)(n \alpha_e(\alpha_{b^{-1}}(m^{-1})), e b^{-1}) \\ = (m, b)(n \alpha_{b^{-1}}(m^{-1}), b^{-1}) \\ = (m \alpha_b(n \alpha_{b^{-1}}(m^{-1})), b b^{-1}) \\ = (m \alpha_b(n) m^{-1}, e) \in \tilde{N} \end{aligned}$$

On the other hand: $(m, b)(e, a)(\alpha_{b^{-1}}(m^{-1}), b^{-1})$
 $= (m, b)(e \alpha_a \alpha_{b^{-1}}(m^{-1}), a b^{-1})$
 $= (m \alpha_b \alpha_a \alpha_{b^{-1}}(m^{-1}), b a b^{-1})$

no reason for this
to be e .

So \tilde{A} is not necessarily normal.

Recognizing semidirect products

We have the notion of semidirect product that is constructed from the data of two groups N, A , and a homomorphism $\alpha: A \rightarrow \text{Aut}(N)$.

We want a way to recognize that a given group is isomorphic to a semidirect product $N \rtimes_{\alpha} A$ for some A, N, α .

Recall: $N \rtimes_{\alpha} A = \{(n, a) \mid n \in N, a \in A\}$ with operation
 $(n, a)(n', a') = (n \alpha_a(n'), aa')$

- There are subgroups $\tilde{N} = \{(n, e) \mid n \in N\}$ $\tilde{A} = \{(e, a) \mid a \in A\}$
- The subgroup \tilde{N} is normal.
- $N \rtimes_{\alpha} A = \tilde{N}\tilde{A}$ and $\tilde{N} \cap \tilde{A} = \{(e, e)\}$ is the trivial subgroup.

The result is that these properties essentially characterize the semidirect product.

Proposition 3.2.5 Let G be a group, and let N and A be subgroups of G .

Suppose (i) N is normal in G

(ii) $NA = G$

(iii) $N \cap A = \{e\}$

Then G is isomorphic to $N \rtimes_{\alpha} A$, where $\alpha: A \rightarrow \text{Aut}(N)$ is the conjugation homomorphism: $\alpha_a(n) = ana^{-1}$.

More specifically, the function

$$\varphi: N \rtimes_{\alpha} A \rightarrow G \quad \varphi(n, a) = na$$

is an isomorphism.

Remark: The hypothesis (i) is necessary in order to know that $c: A \rightarrow \text{Aut}(N)$ makes sense. (If N is not normal then c_a is not necessarily an automorphism of N).

Proof The main thing is to show that $\varphi: N \rtimes_c A \rightarrow G$ is a homomorphism.

$$\begin{aligned} \varphi((n, a)(n', a')) &= \varphi(nc_a(n'), aa') = nc_a(n')aa' \\ &= nan'a^{-1}aa' = nan'a' = \varphi(n, a)\varphi(n', a') \end{aligned}$$

So it works.

The image of φ is clearly NA , which equals G by hypothesis (ii'), so φ is surjective.

Last, $\ker(\varphi) = \{(n, a) \mid na = e\}$. But $na = e$ means $n = a^{-1}$ then $n = a^{-1} \in A$ so $n \in N \cap A = \{e\}$, so $n = e$, and $a = e$ as well. So $\ker(\varphi) = \{(e, e)\}$ is trivial.

Thus φ is injective.

This completes the proof that $\varphi: N \rtimes_c A \rightarrow G$ is an isomorphism. \square