

Primes, modular arithmetic

Proposition 1.6.19 Let $p \in \mathbb{N}$ be prime, let a, b be nonzero integers. If $p \nmid ab$ then $p \nmid a$ or $p \nmid b$.

Proof Since p is prime $\gcd(p, a) = 1$ or p .

if $\gcd(p, a) = p$ then $p \mid a$ and we are done.

if $\gcd(p, a) = 1$, then we can find integers s, t such that

$$1 = ps + at$$

then $b = bps + abt$ (multiply by b)

We know $p \nmid ab$, so $p \nmid abt$, and obviously $p \nmid bps$, so we conclude $p \mid b = bps + abt$. \blacksquare

Theorem 1.6.21 Prime factorization of n is unique
(up to order of the factors).

Proof Strong induction on n . Base case $n=1$: only empty product is possible.

Suppose for every $m < n$, prime factorization is unique.

Suppose $n = p_1 \cdots p_k = q_1 \cdots q_\ell$ are two prime factorizations.

We may reorder the factors so that

$$p_1 \leq p_2 \leq \cdots \leq p_k \quad q_1 \leq q_2 \leq \cdots \leq q_\ell.$$

Suppose $p_1 \leq q_1$ (if not, swap names of p 's and q 's)

Since $p_1 \mid n = q_1 \cdots q_\ell$, we have that $p_1 \mid q_j$ for some $j, 1 \leq j \leq \ell$.
Since p_1 and q_j are prime, $p_1 = q_j$.

Then $p_1 \leq q_1 \leq q_j = p_1$ so $p_1 = q_1 = q_j$.

Then take $m = \frac{n}{p_1} = \frac{n}{q_1} = p_2 \cdots p_k = q_2 \cdots q_l$

since $m < n$, we apply induction hypothesis to conclude $k=l$
and $p_i = q_{l-i}$ for $2 \leq i \leq k$. Thus the two factorizations
of n are not actually different. \square

Modular arithmetic (Clock arithmetic = mod (Z))

Definition let $a, b, n \in \mathbb{Z}$, $n \geq 1$. We say
"a is congruent to b modulo n"
 $a \equiv b \pmod{n}$ if $n | (b-a)$.

Observe: given $a \in \mathbb{Z}$ and $n \geq 1$, we can apply long division
to get $q, r \in \mathbb{Z}$ with $a = qn+r$ and $0 \leq r < n$

Then $a-r = qn$ is divisible by n , so $r \equiv a \pmod{n}$
"Any number a is congruent modulo n to the remainder of
a divided by n."

Integer arithmetic works well with congruence:

Lemma 1.7.5: let $a, a', b, b' \in \mathbb{Z}$. Assume $a \equiv a' \pmod{n}$, $b \equiv b' \pmod{n}$.
then $a+b \equiv a'+b' \pmod{n}$ and $ab \equiv a'b' \pmod{n}$.

Proof Know $n | (a-a')$ and $n | (b-b')$, so
 $n | (a-a') + (b-b') = (a+b) - (a'+b')$
so $(a+b) \equiv (a'+b') \pmod{n}$.

Also $n \mid (a-a')b + a'(b-b')$ $\Rightarrow (ab - a'b + a'b - a'b) = ab - a'b'$
 $\therefore ab \equiv a'b' \pmod{n}$ \blacksquare

Example Compute $(7964) \cdot (11203) \pmod{10}$.

Since $7964 \equiv 4 \pmod{10}$,

$11203 \equiv 3 \pmod{10}$,

we have $7964 \cdot 11203 \equiv 4 \cdot 3 = 12 \equiv 2 \pmod{10}$

Note: $n \mid a \Leftrightarrow a \equiv 0 \pmod{n}$.

Fact: $3 \mid a \Leftrightarrow 3 \mid (\text{sum of digits of } a)$ write $a = \sum_{j=0}^k a_j 10^j$

where $a_j \in \{0, 1, \dots, 9\}$ are digits of a .

Since $10 \equiv 1 \pmod{3}$, we have $10^j \equiv 1^j \equiv 1 \pmod{3}$.

$$\text{So } a = \sum_{j=0}^k a_j 10^j \equiv \sum_{j=0}^k a_j \cdot 1^j \equiv \sum_{j=0}^k a_j \pmod{3}$$

$$\text{so } a \equiv 0 \pmod{3} \Leftrightarrow \sum_{j=0}^k a_j \equiv 0 \pmod{3}.$$

The relation of congruence (i.e. the concept "is congruent to") is the canonical first example of an equivalence relation.

Lemma 1.7.2 For $a, b, c, n \in \mathbb{Z}$, $n \geq 1$, we have

$$(i) \quad a \equiv a \pmod{n} \quad (\text{reflexive})$$

$$(ii) \quad a \equiv b \pmod{n} \text{ iff } b \equiv a \pmod{n} \quad (\text{symmetric})$$

$$(iii) \quad \text{if } a \equiv b \pmod{n} \text{ and } b \equiv c \pmod{n} \text{ then} \\ a \equiv c \pmod{n} \quad (\text{transitive})$$

Proof (i) $a - a = 0$ and $n \mid 0$ (ii) since $b - a = -(a - b)$,

$n \mid b - a \Leftrightarrow n \mid a - b$. (iii) $n \mid b - a$ and $n \mid c - b \Rightarrow n \mid (c - b) + (b - a) = c - a$.

Definition Fix $n \in \mathbb{N}$. For $a \in \mathbb{Z}$, the congruence class of a modulo n is the set

$$[a] = \{ b \in \mathbb{Z} \mid b \equiv a \pmod{n} \} = \{ a + kn \mid k \in \mathbb{Z} \}$$

To emphasize the dependence on n we write $[a]_n$.

Lemma 1.7.3 Fix $n \in \mathbb{Z}$. For $a, b \in \mathbb{Z}$, the following are equivalent.

$$(i) \quad a \equiv b \pmod{n}$$

$$(ii) \quad [a] = [b]$$

$$(iii) \quad \text{rem}_n(a) = \text{rem}_n(b) \quad (\text{rem}_n = \text{remainder upon div. by } n.)$$

$$(iv) \quad [a] \cap [b] \neq \emptyset.$$

Proof Goodman.

Corollary 1.7.4 There are exactly n distinct congruence classes mod n namely $[0], [1], [2], \dots, [n-1]$. These sets are pairwise disjoint.

Denote the set of congruence classes

$$\mathbb{Z}_n = \{[0], [1], \dots, [n-1]\}$$

We wish to define $+$ and \cdot in \mathbb{Z}_n by the formulas

$$[a] + [b] = [a+b] \quad \text{and} \quad [a] \cdot [b] = [ab]$$

This works, but it is worth thinking about why it works. The issue is that the object denoted $[a]$ could be represented other ways, for instance as $[a']$ where $a \equiv a' \pmod{n}$. But then ab would become $a'b'$ which is different....

The crucial point is this: if $[a] = [a']$ and $[b] = [b']$,
 then $[a+b] = [a'+b']$ and $[a \cdot b] = [a' \cdot b']$.

This follows from Lemmas 1.7.3 and 1.7.5.

We say that addition and multiplication of congruence classes is "well-defined". This is a logical pattern that will repeat whenever we try to define a function that seems to depend on a choice of representative element.

Proposition 1.7.7 The operations + and \cdot on \mathbb{Z}_n satisfy
 commutative law, associative law, distributive law.

$[0]$ is additive identity, $[1]$ is multiplicative identity.
 The additive inverse of $[a]$ is $[-a]$ (or $[n-a]$).